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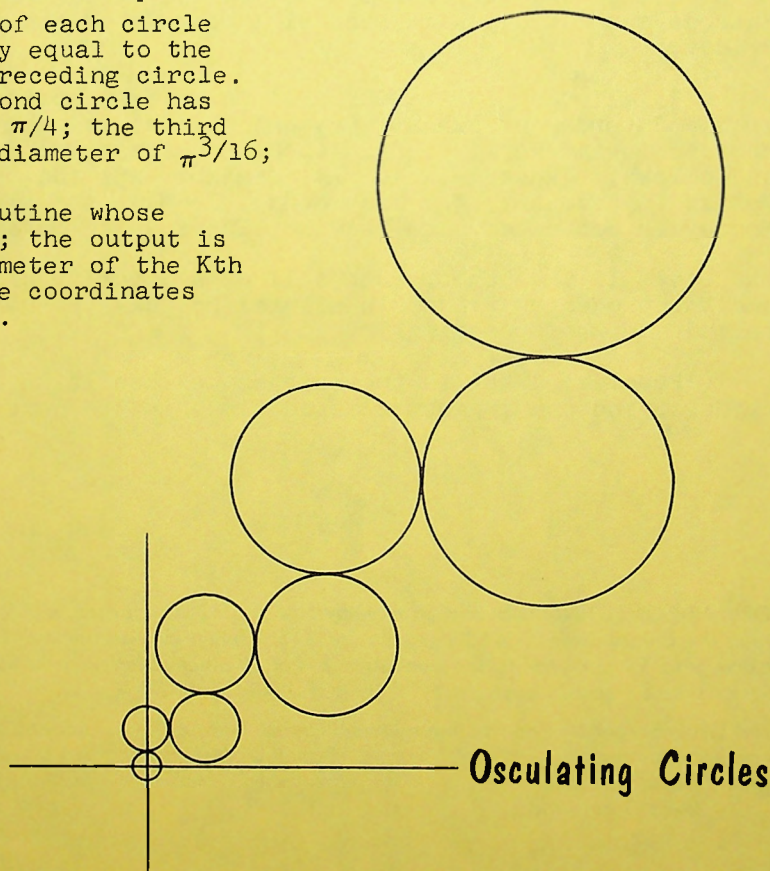
A circle of unit diameter is centered at the origin.

Successive circles are drawn, each tangent to the preceding circle as shown. The points of tangency are alternately on the right and on top.

The diameter of each circle is numerically equal to the area of the preceding circle. Thus, the second circle has a diameter of $\pi/4$; the third circle has a diameter of $\pi^3/16$; and so on.

Write a subroutine whose argument is K; the output is to be the diameter of the Kth circle and the coordinates of its center.

PROBLEM 96



The Payday Problem

A firm pays its employees twice a month, according to the following schedule:

1. Pay checks will normally be issued on the 5th and 20th of each month, if those days are working days.
2. If the 5th or 20th fall on weekends (that is, on Saturdays or Sundays) payment will be made on the preceding Friday.
3. If the 5th or 20th fall on a legal holiday, other than a Saturday or Sunday, payment will be made on the next following working day. Legal holidays are defined as follows: January 1, July 4, December 25, the third Monday in February, the last Monday in May, the first Monday in September, and the last Thursday in November.
4. If the 5th or 20th fall on a Saturday or Sunday, and the preceding Friday is a legal holiday, payment will be made on the preceding Thursday.

Problem: What is the shortest elapsed time between paydays, and the longest elapsed time, in the next 50 years? ☐

PROBLEM 97



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The Altair 8800

PC29-3

The Altair 8800 computer, made by MITS, Inc., P.O. box 8636, 6328 Linn, N.E., Albuquerque, New Mexico 87108, was announced around the first of this year. The basic machine, with 256 8-bit words of storage, sells for \$542 in kit form, or \$755 assembled. The basic machine has volatile storage. Input is by means of toggle switches; output is through binary lights.

The machine we worked with had 1024 words of storage; this size is \$615 in kit form or \$830 assembled. The promotional literature says that 256 words of storage is "enough memory to do many machine language procedures and some control applications." This should be taken rather cautiously; there are few real problems whose solution can be programmed in that size machine.

The operation code length of the Altair is a full word of 8 bits. An instruction, therefore, consists of one or more words. For example, the operation ADD ONE TO ACCUMULATOR is a one-word instruction (octal 074). ADD IMMEDIATE takes two words:

306
037

(add decimal 31 to the contents of the accumulator). Most of the normal instructions that one expects on a computer take three words, such as:

072
126
002

for which the 072 is LOAD ACCUMULATOR, and the other two words indicate address 1126 octal, as follows:


00	000	010		01	010	110
0	0	2		1	2	6

(In the basic machine with 256 words of storage, the third of those three instruction words would always be zero.)

The 8800 uses the Intel 8080 chip as its central processor, and hence the Altair uses the op-codes of the Intel chip. Many of these are of marginal use in a general-purpose machine, such as RETURN (from subroutine) ON ZERO, RETURN ON MINUS, RETURN ON PLUS, and the like. Simple instructions like ADD take four words:

041
123
001
206

(add contents of word 523 to the accumulator).



The machine has a HALT instruction (166), but on execution it tends to lock the entire machine with all console lights on and all switches inoperative. This condition can be relieved by hitting STOP and RESET simultaneously. For purposes of halting execution, a better combination is:

A	303
A+1	A
A+2	---

(unconditional jump to current location) which ties the machine into a short loop, during which the manual STOP switch can be used.

The machine is quite reliable, but its circulating storage is not protected by parity checks, and the loss or gain of a single bit can be disastrous.

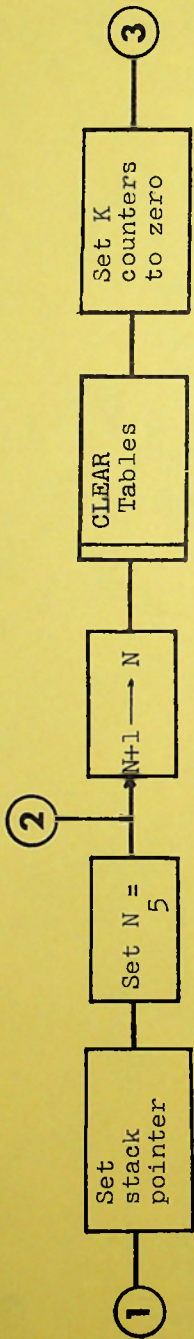
Promotion for the machine has leaned toward over-enthusiasm and hyperbole. For example, "It has 78 basic machine instructions with variances up to 200 instructions. That's enough to program all the traffic lights in a major city." The relation between the number of op-codes and any specific application is obscure.

Or, "Number of subroutine levels: 65,000." To call a subroutine takes 3 words; to return from a subroutine takes one word, and the return address (two words) must be stored somewhere. Thus, even with a full size machine with 65,000 words, it is difficult to determine just what the advertised statement could mean.

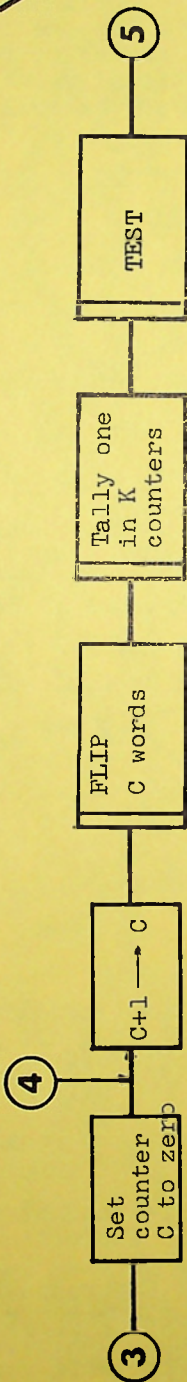
Additionally, the reference manual is in need of extensive re-writing; the explanations of the actions for the various op-codes are poorly written and hard to decipher even by experienced programmers.

(from the manual) "...the Altair 8800 can execute a six instruction addition program approximately 30,000 times in one second." This is true, but implies that the operating speed is 180,000 instructions per second; for a typical program, it will be more like 116,000 instructions per second.

Again, "Zero bit--this bit is set to 1 if the result of certain instructions is zero..." (underscores mine). The user of a computer wants to know the precise action of every op-code, and those things that can affect any bit.



A

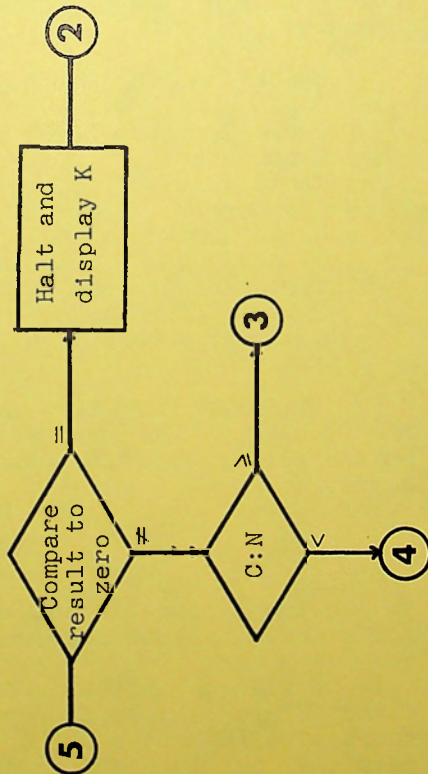


MAIN flowchart for Penny Flipping I.

Note: the bar at the left end of a rectangle denotes "go to a subroutine."

N is the chief parameter of the problem; i.e., the size of the stack of pennies. The tables are 80 words each, allowing N to go to 80.

The TEST subroutine returns 1 in a trigger for failure; 0 for success (that is, the stack returns to all heads).



The original Penny Flipping problem seems ideally suited to the machine. Given a stack of N pennies, all heads up. The top penny is turned over, then the top two pennies, the top three, and so on until the entire stack is turned over, after which the top one is turned over, the top two, and so on. The number of flips is counted, and the result is the functional value for that N. The results for various values of N are as follows:

N	f	N	f
1	2	11	121
2	3	12	119
3	9	13	116
4	11	14	195
5	24	15	75
6	35	16	79
7	28	17	204
8	31	18	323
9	80	19	228
10	60	20	199

Flowcharts are given for this problem, to extend the function from 21 to 80. Since the counts of the number of flips can go over 127 (the limit of the size of a positive number in one word), a two-word counter is set up, and subroutine C shows the logic of incrementing those counters (the scheme readily extends to larger counters).

We begin by allocating storage as follows:

Table T, words 020 through 137 octal (this table will contain the stack of pennies, one penny per word).

Table TC (T copy), words 140 through 257 octal.

K counters, words 002 and 003.

N, word 004.

C, word 005.

Temporary storage, word 006.

Trigger (for TEST), word 010.

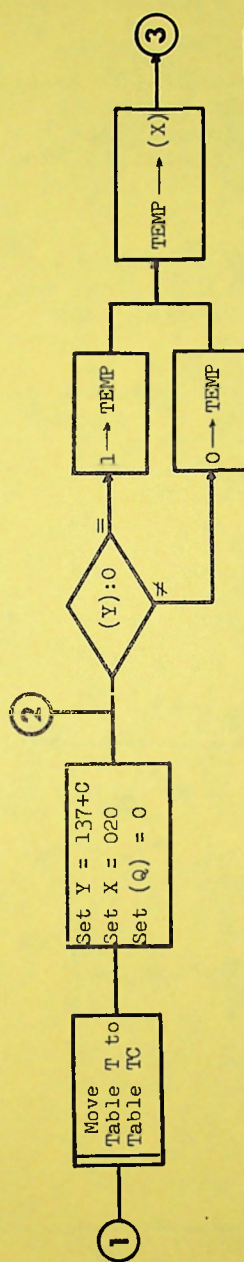
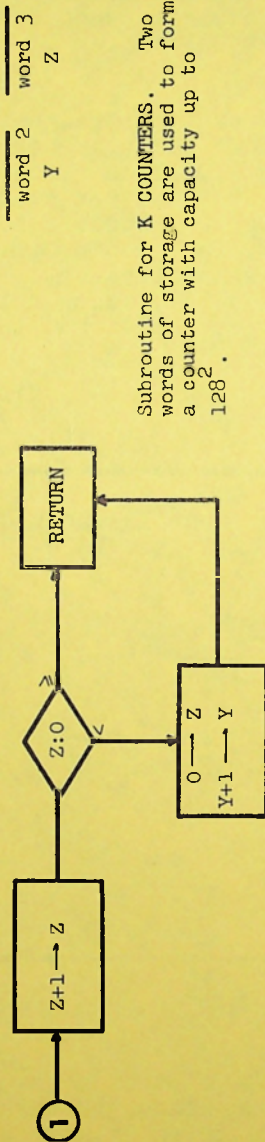
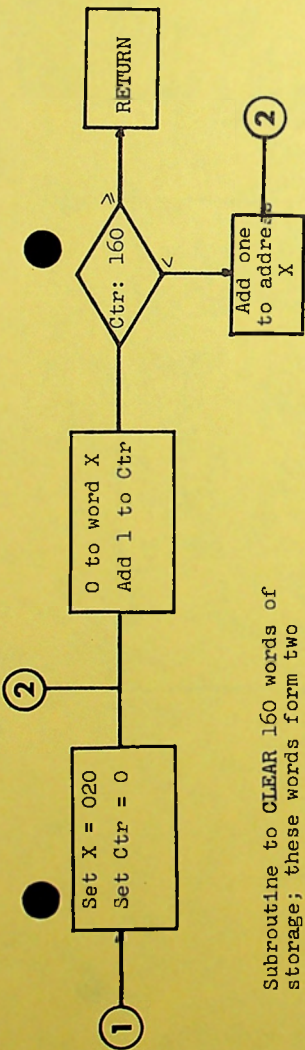
Counter, word 011.

The MAIN flowchart (A) shows the overall logic of the problem. The program begins (as must all programs on the Altair that use subroutines) by setting the stack pointer:

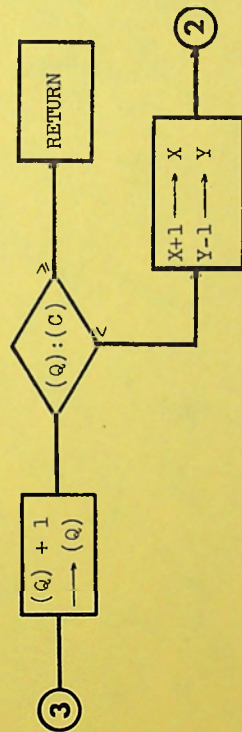
```
061
377
003
```

to provide a place in storage (in this case at octal 1777) for stacking return addresses for all subroutines.

N will be initialized to 6, at Reference 2. Tables T and TC are cleared to zero. The flowchart for this subroutine (B) shows a loop that counts to 160. This is possible, but it might be easier (if this is your first approach to the machine) to form two loops, each counting to 80)₁₀.



Subroutine to FLIP C words of Table T.
(Parentheses indicate "contents of")



PC29-8

Subroutine (D) (FLIP) performs the heart of the problem. It begins by calling on another subroutine (for which the flowchart is not given) to move table T to TC. Then a loop is initialized to the first address of Table T (020) and the last address of Table TC ($137 + C$), to flip C words. The counter Q is word 011.

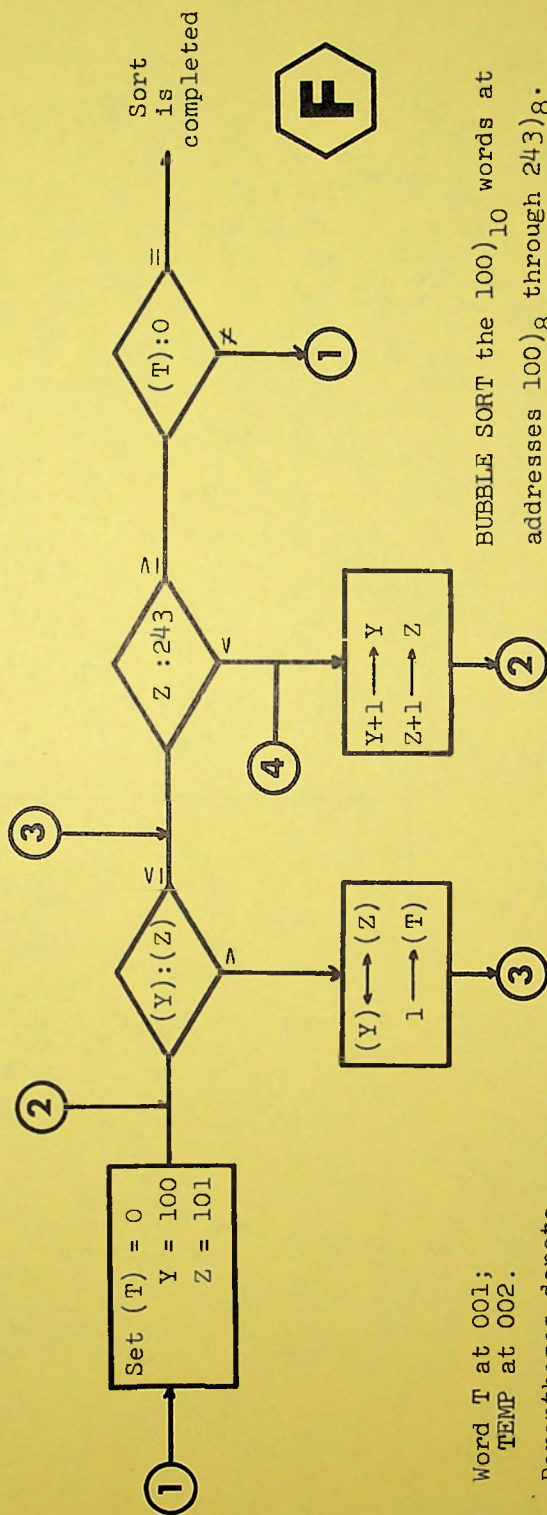
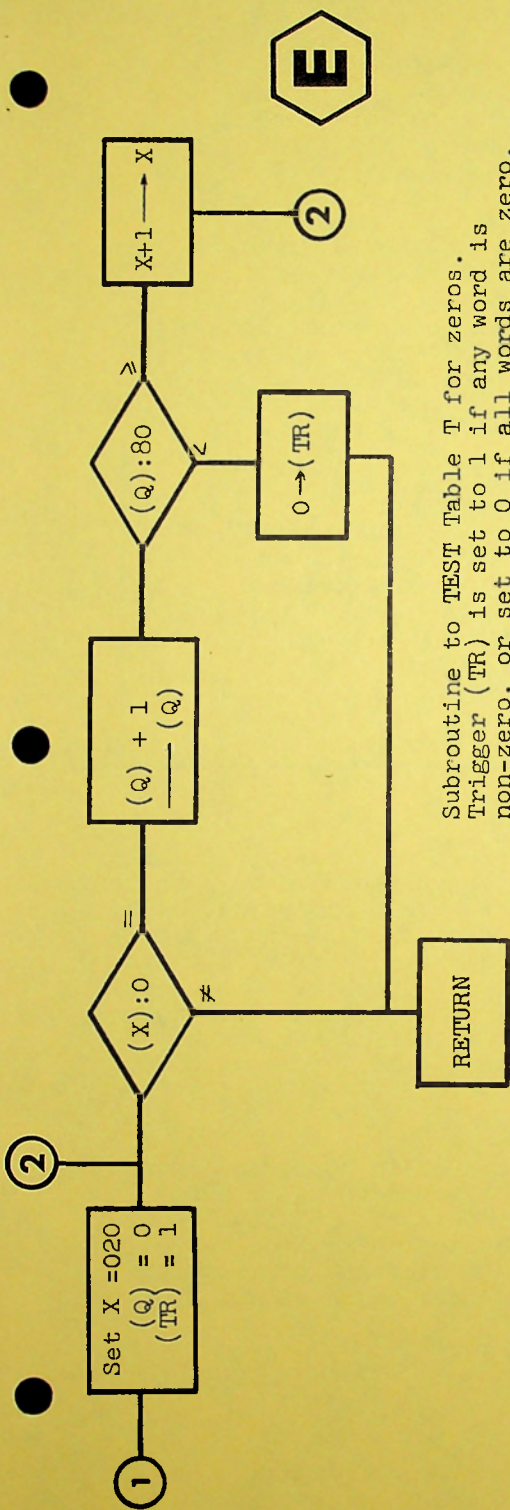
The TEST subroutine (E) examines all of Table T for zeros. The flipping task is completed when N words of the table are again zero, but it is simpler to test the entire table.

The complete program, written in straightforward fashion (that is, not using tricky coding) occupies about 400 words of storage, plus some 170 words of data. It executes at an average rate of 80 flips per second, so that the total run for $N = 74$ is 70 seconds (the result is 5475). For each value of N, the machine is halted, and the contents of the counters is displayed. For $N = 74$, this will read as 052 and 143; these results in octal are changed to 42 and 99 (decimal), and the result is then $42 \cdot 128 + 99 = 5475$.

Further results have also been obtained for the Penny Flipping IV problem (see PC25-12) as follows:

	N	f	N	f	
Allan Herschderfer and Charles McCord, April 1975, on a CDC 3170	38	163540	45	88685	FJG, May 1975, 8800 on the Altair
	39	20064	46	4860	
	40	31122	47	15088	
	41	51040	48	518880	
	42	11480	49	43776	
	43	16128	50	25284	
	44	450296	51	82000	
			52	5100	
			53	73008	
			54	2270520	
			55	517083	

For another example of the level that most users of the Altair 8800 will work at, consider the task of sorting one hundred numbers. The given data is stored in words $100)_8$ through $243)_8$, and we will produce a machine language code to bubble sort it.



Bubble (or interchange) sorting is considered by many (including Knuth) to be a dirty technique that should not be encouraged in any way. Granted that bubble sorting is crude, inelegant, unsophisticated, and (for more than a few items to be sorted) highly wasteful of machine time, it still exists and has definite virtues. For one, it is simple and easy to code in any language, and hence easy to debug and test. Its inherent inefficiency hurts you only if the number of things to be sorted is large (say, over 30) or the number of sets of those things is large (say, over 1000), or both. Otherwise, the programmer is choosing to swap machine time (and possibly storage space) for brain strain, or for a reduction in elapsed time--and that swap can be profitable. Finally, the technique is widely used, and there should be no hidden or suppressed tricks in our trade.

The accompanying flowchart (F) shows the logic of the bubble sort, together with a complete code for it (G). The code, like the technique itself, is simple, using no esoteric coding tricks, nor any of the hidden registers of the Altair. If one chooses to use an inefficient technique, then it is inconsistent to try to use it in a sophisticated way.

Finally, a problem is offered that is particularly suited to the Altair 8800, but is solvable on any machine. The accompanying diagram (H) shows 32 8-bit words. The words are indexed at the left, from 00 through 37 octal. The individual bits in the words are also indexed at the top, from 000 through 111 binary. The heavy line divides each word into two parts. The 5-bit part is the address of one of the 32 words; the 3-bit part addresses one of the bits of that word. For example, the word at 00000 addresses the 5th bit from the left in word 10011 (that bit is circled in the diagram).

The 32 words are to be scanned in turn, and each indicated bit is to be changed, from 1 to 0, or from 0 to 1. The implementation of this simple situation is a nice exercise in bit manipulation. The Problem is: Will the process converge for any starting combination of bits (other than the obvious case in which all 256 bits are initially zero)?

500	076	R-1
501	000	$\text{Set } (T) = 0$
502	062	
503	001	
504	000	
505	076	$\text{Set } Y = 100$
506	100	
507	062	
510	134	
511	001	
512	062	
513	151	
514	001	
515	062	$\text{Set } Z = 101$
516	162	
517	001	
520	076	
521	101	$\text{Set } Z = 101$
522	062	
523	137	
524	001	
525	062	
526	157	
527	001	
530	062	
531	170	$\text{Set } Z = 101$
532	001	
533	072	
534	000	
535	000	$R-2$
536	041	
537	000	
540	000	
541	226	$(Y) - (Z)$
542	312	
543	177	
544	001	
545	372	$\text{To } R-3$ on zero or minus
546	177	
547	001	
550	072	
551	000	$(Y) \rightarrow (\text{TEMP})$
552	000	
553	062	
554	002	
555	000	$(Z) \rightarrow (Y)$
556	072	
557	000	
560	000	
561	062	$(Z) \rightarrow (Y)$
562	000	
563	000	

564	072	$(\text{TEMP}) \rightarrow (Z)$
565	002	
566	000	
567	062	
570	000	$1 \rightarrow (T)$
571	000	
572	076	
573	001	
574	062	$R-5$
575	001	
576	000	
577	072	
600	137	$Z: 243$
601	001	
602	326	
603	243	
604	372	$G_0 R-4$
605	222	
606	001	
607	072	
610	001	$(T): 0$
611	000	
612	306	
613	000	
614	312	$G_0 R-1$
615	257	
616	001	
617	303	
620	100	$R-4$
621	001	
622	072	
623	134	
624	001	$Y+1 \rightarrow Y$
625	074	
626	062	
627	134	
630	001	$Y+1 \rightarrow Y$
631	062	
632	151	
633	001	
634	062	$Y+1 \rightarrow Y$
635	162	
636	001	

637	072	$Z+1 \rightarrow Z$
640	137	
641	001	
642	074	
643	062	$Z+1 \rightarrow Z$
644	137	
645	001	
646	062	
647	157	$G-R-2$
650	001	
651	062	
652	170	
653	001	$G-R-2$
654	303	
655	133	
656	001	
657	Continue	



Code for the Altair 8800
to bubble sort the words
from 100-243)8 in
ascending order.

	000	001	010	011	100	101	110	111
00000	1	0	0	1	0	0	1	1
00001	1	0	1	1	1	1	1	0
00010	1	1	0	0	0	0	0	0
00011	1	0	0	1	1	1	1	1
00100	0	0	1	1	1	0	0	0
00101	1	1	0	0	1	1	1	1
00110	1	1	0	0	1	1	0	0
00111	1	0	0	0	0	0	1	1
01000	1	1	1	0	0	0	0	1
01001	1	0	0	0	1	1	1	0
01010	1	1	0	0	1	0	0	1
01011	1	1	0	1	0	1	1	0
01100	1	0	1	0	1	0	0	0
01101	1	1	1	0	0	1	0	1
01110	0	1	1	0	1	1	0	1
01111	0	1	1	0	0	0	0	1
10000	0	1	0	1	0	1	1	0
10001	1	1	1	1	1	1	0	1
10010	0	1	0	0	0	0	1	1
10011	1	0	1	0	1	0	1	0
10100	0	0	1	0	1	0	0	0
10101	0	1	0	1	1	0	0	1
10110	0	0	0	1	1	0	0	1
10111	1	1	0	1	1	0	1	1
11000	1	0	0	1	0	1	1	0
11001	1	1	0	1	0	1	0	0
11010	0	1	1	0	1	1	0	0
11011	0	0	1	1	1	1	1	1
11100	0	0	0	1	0	0	0	1
11101	0	1	0	0	0	1	0	0
11110	0	1	1	0	0	0	1	0
11111	1	1	1	1	0	1	0	1



NOTONE Research

PC29-13

The game of NOTONE was proposed by Walter Koetke in the first issue (Nov/Dec 1974) of Creative Computing. Two players take turns tossing two dice. When it is a player's turn, he tosses the dice and establishes his point. He may then toss the dice as many times as he wishes, and his score for the turn is the sum of the tosses. If his point reappears, however, his turn ends then with a score of zero. A game consists of ten rounds for each player.

A program (in BASIC) for the game was presented in the Mar/Apr 1975 issue, written by Robert Puopolo. The printout showed a score of 345 for the human player, indicating an average score of 34.5 per turn.

What is the proper strategy of play for this game?

PROBLEM 99

If the point is 7, there is a .50 probability of making a run of 3 tosses before a 7 reappears. If the point is 2 (or 12) however, one can expect a run of 24 tosses before the reappearance of the 2 (or 12), again with a .50 probability. These probabilities are summarized in the accompanying table (A). For example, with a point of 5 (or 9), and a .45 probability, a run of 6 tosses can be expected. Thus, the attempted run at each turn is controlled by the point that is first made, as well as the probability level that will maximize the sum. Notice that although the probability of a run of a given length is the same for points 3 and 11, the expected sum is not. For example, if a .60 probability is used, the expected run for 3 or 11 is 9 tosses, but the 9 tosses starting with 11 will total higher than the 9 tosses starting with 3.

The point about dice is that they turn two flat distributions (that is, two distributions uniformly distributed in the range from 1 to 6) into a peaked distribution from 2 to 12 with these theoretical frequencies:

2	3	4	5	6	7	8	9	10	11	12
1/36	2/36	3/36	4/36	5/36	6/36	5/36	4/36	3/36	2/36	1/36

B

2	7	10	13	18	24	28	32	40
3	4	5	6	9	12	14	16	20
4	2	3	4	6	7	9	10	14
5	1	2	3	5	5	6	7	10
6	1	1	2	4	4	5	6	8
7	1	1	1	3	3	4	5	7

A

Most problems involving dice tosses (those that do not yield to an analytic solution) are best explored by computer. At times, however, some research can be done by hand, given access to suitable dice tosses. The accompanying table (B) gives 1800 tosses of two dice, for which the observed frequencies are as follows:

2	3	4	5	6	7	8	9	10	11	12
37	98	132	200	247	302	268	216	154	98	48

(the chi-squared value for goodness-of-fit is 8.432.)

In the traditional game of craps, the spots appearing on two dice are added. If they were multiplied, the tosses would take this distribution:

1	1/36	8	2/36	18	2/36
2	2/36	9	1/36	20	2/36
3	2/36	10	2/36	24	2/36
4	3/36	12	4/36	25	1/36
5	2/36	15	2/36	30	2/36
6	4/36	16	1/36	36	1/36

Suppose, then, we define the game of CRAPS2 as follows:

The player with the dice makes a toss. If he gets a square (1, 4, 9, 16, 25, or 36) he loses but retains the dice. If he throws 6 or 12, he wins. In all other cases, he has established his point and continues to throw. If he makes his point before the appearance of either 6 or 12, he wins; otherwise he loses.

The usual dice questions are then:

1. What are the odds against winning?
2. What are the odds against making the points of 2, 3, 5, 8, 10, 15, 18, 20, 24, and 30?
3. What are the expected lengths of runs between successive appearances of each of the 18 possible numbers?

Traditionally, dice are used by twos. Suppose they were tossed three at a time, with the spots added? The possibilities then range from 3 to 18 with these frequencies:

3	1/216	8	21/216	13	21/216
4	3/216	9	25/216	14	15/216
5	6/216	10	27/216	15	10/216
6	10/216	11	27/216	16	6/216
7	15/216	12	25/216	17	3/216
		18	1/216		

The same questions (above) could then be raised for CRAPS3, played like normal craps but with these variations: On the first toss, totals of 10 or 11 are an immediate win; totals of 3 or 18 are an immediate loss. Otherwise, a point has been established, and the player keeps tossing until he makes his point (a win) or until 10 or 11 appears (a loss).

Finally, the game of NOTONE can be played with two (multiplicative) dice, or with three (additive) dice, and the proper strategy for these two games is to be determined.

* * * * *

A computer run was made on tosses of two dice, to determine empirically the length of strings between successive appearances of each possible number. The following table shows the statistics gathered on 36000 tosses:

Point	N	Sum	Sum of squares
2	975	27278	1059380
3	2006	33753	984789
4	3000	35543	766037
5	4057	35850	598062
6	4984	35985	474381
7	5968	35978	393352
8	5094	36025	469881
9	3969	35907	615531
10	2999	35583	777941
11	1956	33940	1003532
12	992	27396	1055050

For example, the point 2 appeared 975 times; the sum of the lengths of the strings of tosses between successive 2's was 27278 and the sum of the squares of those string lengths was 1059380. Thus, the average run length between successive 2's was 27.977 and the standard deviation was 26.93.

Combinatorial

PC29-17

There are 63 numbers that are made up of the factors 2, 3, 5, 7, 11, and 13, each taken no more than once. For example, eight of the 63 numbers are:

$7 = 7$
 $39 = 3 \times 13$
 $143 = 11 \times 13$
 $154 = 2 \times 7 \times 11$
 $1365 = 3 \times 5 \times 7 \times 13$
 $2002 = 2 \times 7 \times 11 \times 13$
 $15015 = 3 \times 5 \times 7 \times 11 \times 13$
 $30030 = 2 \times 3 \times 5 \times 7 \times 11 \times 13$

These 63 numbers total 96767; their mean is 1535.9841 and their standard deviation is 4210.7.

(1) What will be the sum, mean, and standard deviation for the 63 numbers similarly made up of the factors 3, 5, 7, 11, 13, and 17?

(2) Write a Fortran program that will utilize six factors (I, J, K, L, M, and N, all integers, all relatively prime to each other) and calculate the sum, mean, and standard deviation of the 63 numbers that can be formed.

PROBLEM 103

* * * * *

N-SERIES 29

Log 29 1.462397997898956087332846762969254991254294417887154
 Ln 29 3.367295829986474027183272032361911605494512913922744
 $\sqrt{29}$ 5.385164807134504031250710491540329556295120161644789
 $\sqrt[3]{29}$ 3.072316825685847293312637982105597485502783238876096
 $\sqrt[4]{29}$ 1.961009057454548013206356850097824143767813931689090
 $\sqrt[10]{29}$ 1.400360331291395876495106634543574736747598468978608
 $\sqrt[100]{29}$ 1.034246309741025914720034059677913189875415336590493
 e^{29} 3931334297144.042074388620580843527685796942333443902
 185643090885541997881948228824965987972
 π^{29} 261424513284460.8709628724332156234476313094693009254
 5656390579769711436771704172403172432
 $\tan^{-1} 29$ 1.536327225795388605121153267285351612646716182990619

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